

NoSQL & NewSQL



NoSQL

Performance Comparison

On 50+ GB data:

- MySQL
 - Writes 300 ms avg
 - Reads 350 ms avg
- Cassandra
 - Writes 0.12 ms avg
 - Reads 15 ms avg

We Don't Want No SQL !

- NoSQL movement: SQL considered slow → only access by id („lookup“)
 - Deliberately abandoning relational world: „too complex“, „not scalable“
 - No clear definition, wide range of systems
 - Values considered black boxes (documents, images, ...)
 - simple operations (ex: key/value storage), horizontal scalability for those
 - ACID → CAP, „eventual consistency“

- Systems

-
- The diagram illustrates the categorization of NoSQL systems based on their data model. It features three yellow boxes at the top labeled 'documents', 'columns', and 'key/values'. Below each box is a red-bordered box containing specific database systems. 'documents' points to 'MongoDB, CouchDB'; 'columns' points to 'Cassandra, HBase'; and 'key/values' points to 'Riak, Redis'. A separate bullet point lists 'Proprietary: Amazon, Oracle, Google, Oracle NoSQL'.
- Open source MongoDB, CouchDB, Cassandra, HBase, Riak, Redis
 - Proprietary: Amazon, Oracle, Google, Oracle NoSQL

- See also: <http://glennas.wordpress.com/2011/03/11/introduction-to-nosql-john-nunemaker-presentation-from-june-2010/>

NoSQL

- Previous „young radicals“ approaches subsumed under „NoSQL“
- = we want „no SQL“
- Well..., „not only SQL“
 - After all, a QL is quite handy
 - So, QLs coming into play again (and 2-phase commits = ACID!)
- Ex: MongoDB: „tuple“ = JSON structure

```
db.inventory.find(  
  { type: 'food',  
    $or: [ { qty: { $gt: 100 } }, { price: { $lt: 9.95 } } ]  
  } )
```

Another View: Structural Variety in Big Data

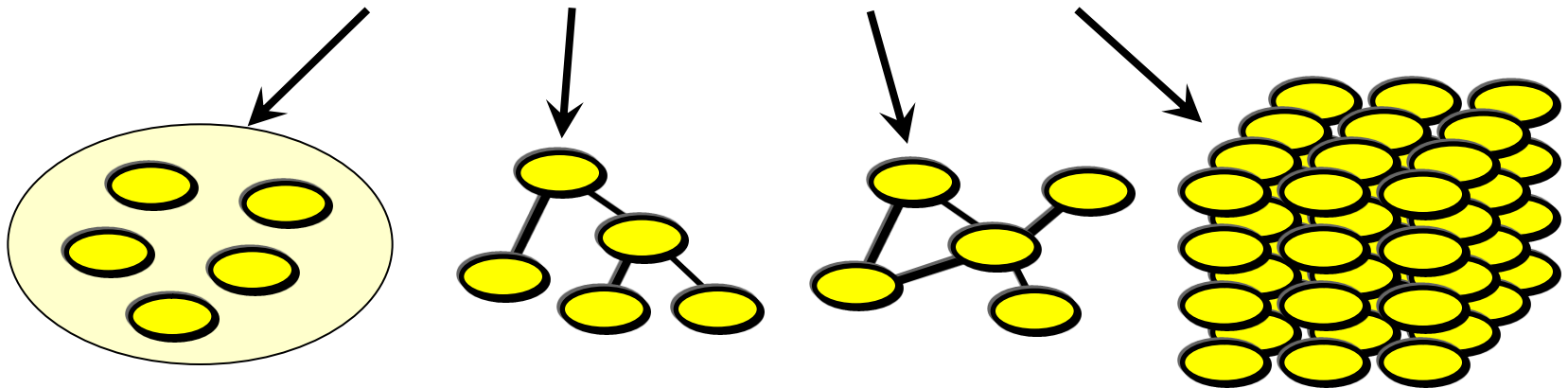
- Stock trading: 1-D sequences (i.e., **arrays**)
- Social networks: large, homogeneous **graphs**
- Ontologies: small, heterogeneous **graphs**
- Climate modelling: 4D/5D **arrays**
- Satellite imagery: 2D/3D **arrays** (+irregularity)
- Genome: long string **arrays**
- Particle physics: **sets** of events
- Bio taxonomies: **hierarchies** (such as XML)
- Documents: key/value stores = **sets** of unique identifiers + whatever
- etc.

Another View: Structural Variety in Big Data

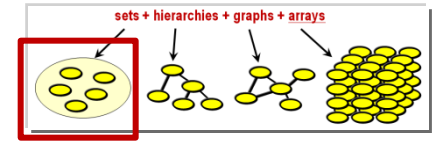
- Stock trading: 1-D sequences (i.e., **arrays**)
- Social networks: large, homogeneous **graphs**
- Ontologies: small, heterogeneous **graphs**
- Climate modelling: 4D/5D **arrays**
- Satellite imagery: 2D/3D **arrays** (+irregularity)
- Genome: long string **arrays**
- Particle physics: **sets** of events
- Bio taxonomies: **hierarchies** (such as XML)
- Documents: key/value stores = **sets** of unique identifiers + whatever
- etc.

Structural Variety in [Big] Data

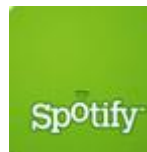
sets + hierarchies + graphs + arrays



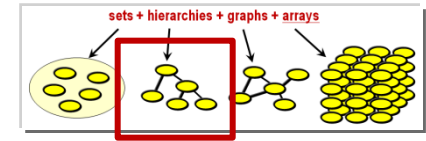
Ex 1: Key/Value Store



- Conceptual model: key/value store = **set of key+value**
 - Operations: *Put(key,value)*, *value = Get(key)*
 - → large, distributed **hash table**
- Needed for:
 - twitter.com: tweet id -> information about tweet
 - kayak.com: Flight number -> information about flight, e.g., availability
 - amazon.com: item number -> information about it
- Ex: Cassandra (Facebook; open source)
 - Myriads of users, like:



Ex 2: Document Stores

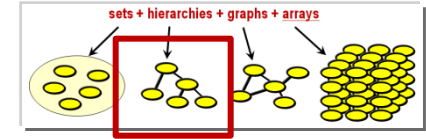


- Like key/value, but value is a **complex document**
 - Data model: set of nested records
- Added: **Search** functionality within document
 - Full-text search: Lucene/Solr, ElasticSearch, ...
- Application: content-oriented applications
 - Facebook, Amazon, ...
- Ex: MongoDB, CouchDB

```
db.inventory.find( { $and: [ { status: "A" }, { qty: { $lt: 30 } } ] } )
```

```
SELECT * FROM inventory WHERE status = "A" AND qty < 30
```

Ex 3: Hierarchical Data



- Disclaimer: long before NoSQL!

```
doc ("books.xml") /bookstore/book/title
```

```
doc ("books.xml") /bookstore/book[price<30]
```

- Later more, time permitting!

```
<?xml version="1.0" encoding="UTF-8"?>

<bookstore>

<book category="COOKING">
  <title lang="en">Everyday Italian</title>
  <author>Giada De Laurentiis</author>
  <year>2005</year>
  <price>30.00</price>
</book>

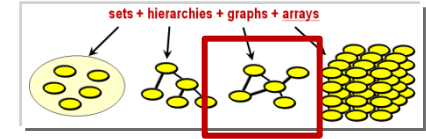
<book category="CHILDREN">
  <title lang="en">Harry Potter</title>
  <author>J K. Rowling</author>
  <year>2005</year>
  <price>29.99</price>
</book>

<book category="WEB">
  <title lang="en">XQuery Kick Start</title>
  <author>James McGovern</author>
  <author>Per Bothner</author>
  <author>Kurt Cagle</author>
  <author>James Linn</author>
  <author>Vaidyanathan Nagarajan</author>
  <year>2003</year>
  <price>49.99</price>
</book>

<book category="WEB">
  <title lang="en">Learning XML</title>
  <author>Erik T. Ray</author>
  <year>2003</year>
  <price>39.95</price>
</book>

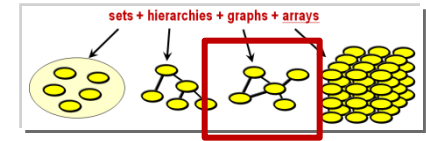
</bookstore>
```

Ex 4: Graph Store

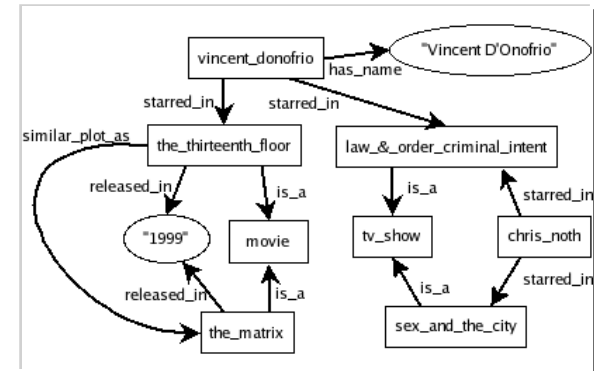


- Conceptual model: Labeled, directed, attributed graph
- Why not relational DB? can model graphs!
 - but “endpoints of an edge” already requires join
 - No support for global ops like transitive hull
- Main cases:
 - Small, heterogeneous graphs
 - Large, homogeneous graphs

Ex 4a: RDF & SPARQL



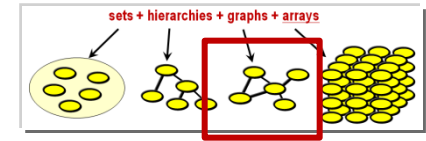
- Situation: **Small, heterogeneous** graphs
- Use cases: ontologies, knowledge graphs, Semantic Web
- Model:
 - Data model: graphs as triples
→ RDF (Resource Data Framework)
 - Query model: patterns on triples
→ SPARQL (see later, time permitting)



```

PREFIX foaf: <http://xmlns.com/foaf/0.1/>
SELECT ?name ?mbox
WHERE
{
    ?x foaf:name ?name .
    ?x foaf:mbox ?mbox
}
  
```

Ex 4b: Graph Databases

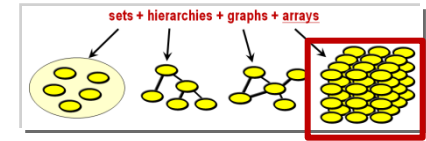


- Situation: **Large, homogeneous** graphs
- Use cases: Social Networks
- Common queries:
 - *My friends*
 - *who has no / many followers*
 - *closed communities*
 - *new agglomerations,*
 - *new themes, ...*
- Sample system: Neo4j with QL Cypher



```
MATCH (:Person {name: 'Jennifer'})-[:WORKS_FOR]->(company:Company)
RETURN company.name
```

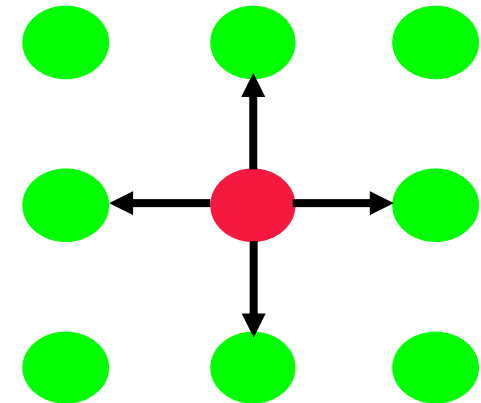
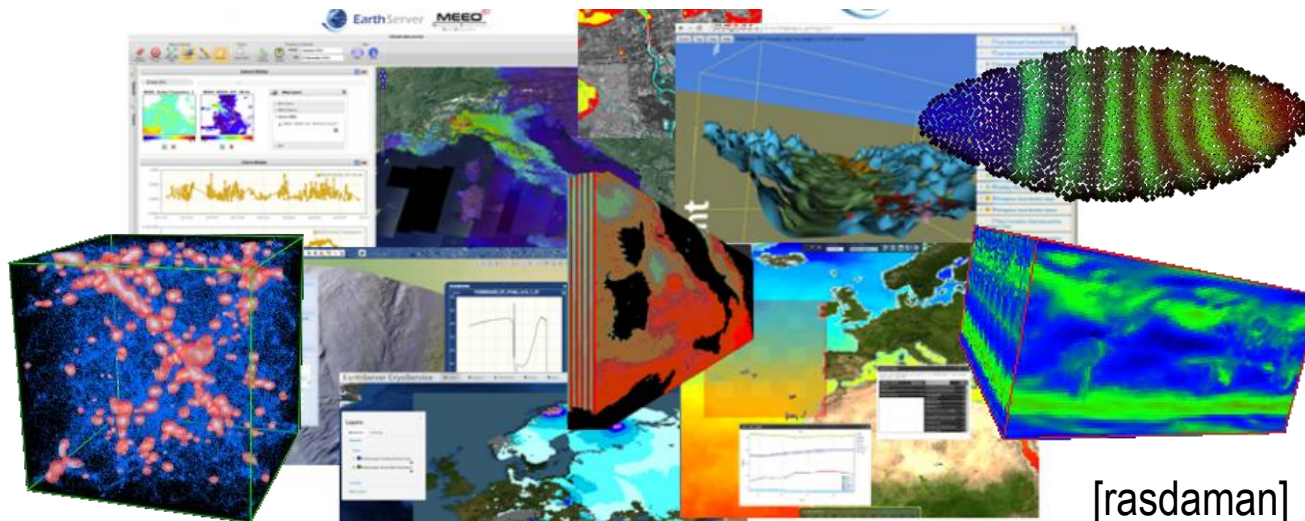
Ex 5: Array Analytics



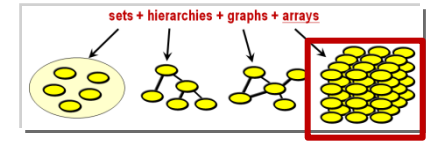
- **Array Analytics** :=
Efficient analysis on multi-dimensional arrays of a size several orders of magnitude above the evaluation engine's main memory

sensor, image [timeseries],
simulation, statistics data

- Essential property: n -D Cartesian neighborhood



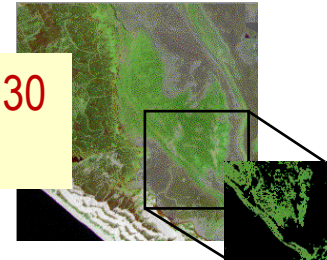
Ex 5: Array Databases



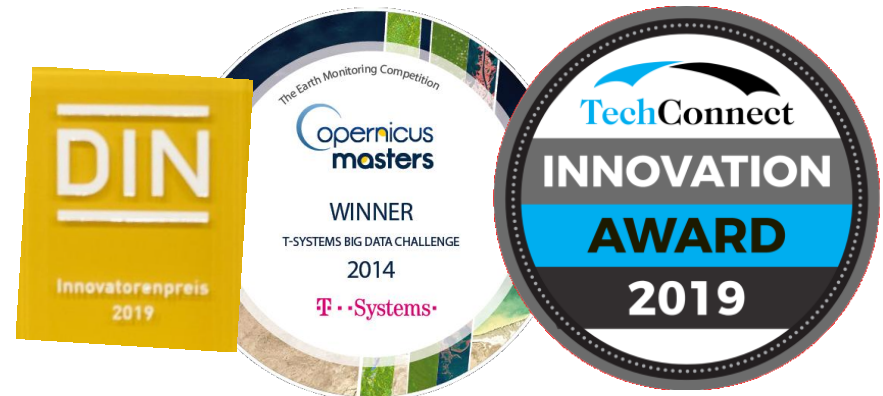
- Ex: rasdaman = **Array DBMS**

- Data model: n-D arrays as attributes
- Query model: Tensor Algebra
- Demo at <http://standards.rasdaman.org>

```
select img.raster[x0:x1,y0:y1] > 130
from LandsatArchive as img
```



- Multi-core, distributed, platform for EarthServer (<https://earthserver.xyz>)
- Relational? „Array DBMSs can be 200x RDBMS“ [Cudre-Maroux]



Arrays in SQL



- 2014 - 2018
- rasdaman as blueprint

```
create table LandsatScenes(
  id: integer not null, acquired: date,
  scene: row( band1: integer, ..., band7: integer ) mdarray [ 0:4999,0:4999] )
```

```
select id, encode(scene.band1-scene.band2)/(scene.band1+scene.band2)), „image/tiff“ )
from LandsatScenes
where acquired between „1990-06-01“ and „1990-06-30“ and
      avg( scene.band3-scene.band4)/(scene.band3+scene.band4)) > 0
```

NewSQL

NewSQL: *The Empire Strikes Back*

- Michael Stonebraker: „*no one size fits all*“
- NoSQL: sacrificing functionality for performance – no QL, only key access
 - Single round trip fast, complex real-world problems slow
- Swinging back from NoSQL:
declarative QLs considered good (again), but SQL often inadequate
- Definition 1: NewSQL = SQL with **enhanced performance** architectures
- Definition 2: NewSQL = SQL enhanced with, eg, **new data types**
 - Some call this NoSQL

What Makes an RDBMS Slow?

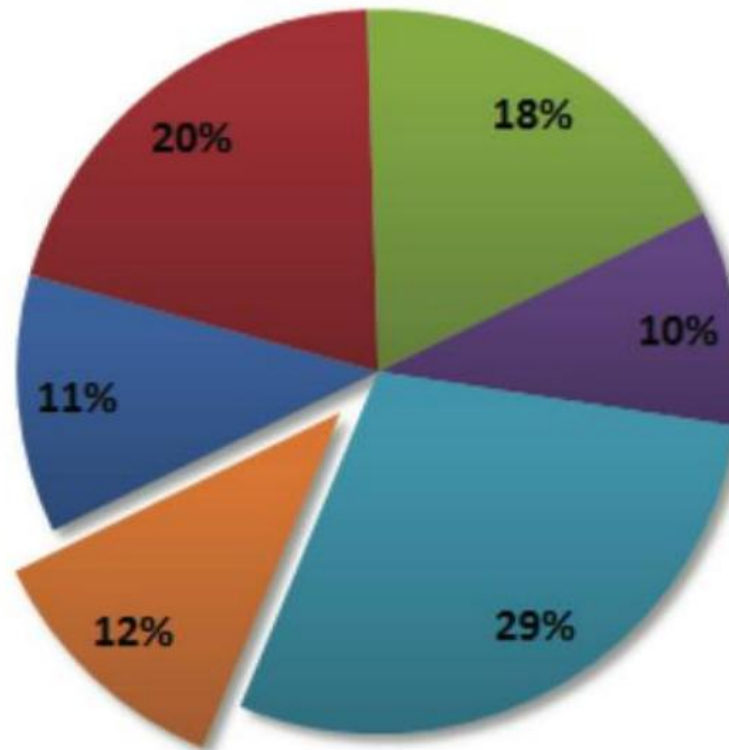
General Purpose RDBMS Processing Profile

OLTP Through the Looking Glass, and What We Found There

Stavros Harizopoulos, Daniel Abadi, Samuel Madden, and Michael Stonebraker

ACM SIGMOD 2008.

- Index Management
- Logging
- Locking
- Latching
- Buffer Management
- Useful Work

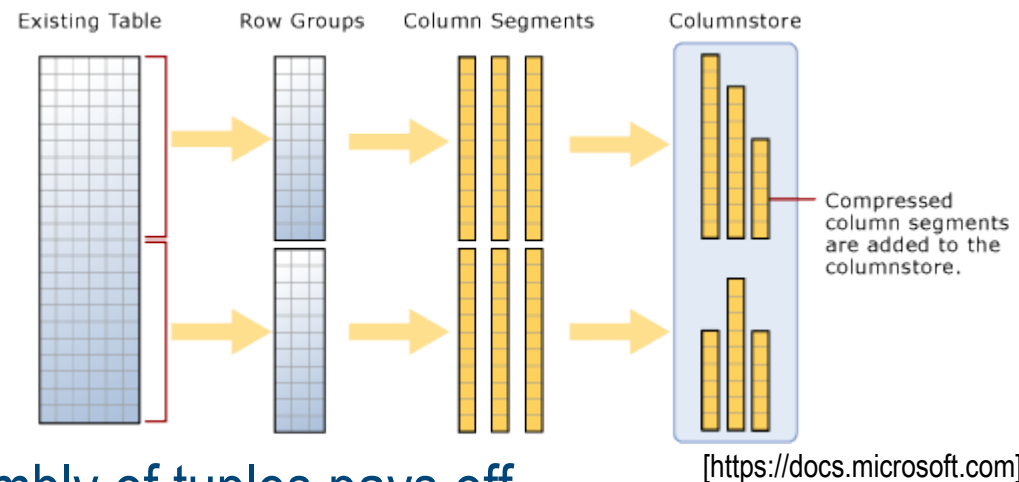


Column-Store Databases

- Observation: fetching long tuples overhead when few attributes needed
- Brute-force decomposition: one value (plus key)

- Ex: Id+SNLRH → Id+S, Id+N, Id+L, Id+R, Id+H

- Column-oriented storage:
each binary table separate file



- With clever architecture, reassembly of tuples pays off
 - system keys, contiguous, not materialized, compression, MMIO, ...

- Sample systems: MonetDB, Vertica, SAP HANA

Main-Memory Databases

- RAM faster than disk → load data into RAM, process there
 - CPU, GPU, ...
- Largely giving up ACID's Durability → different approaches
- Sample systems: ArangoDB, HSQLDB, MonetDB, SAP HANA, VoltDB, ...

Database Landscape Map – December 2012

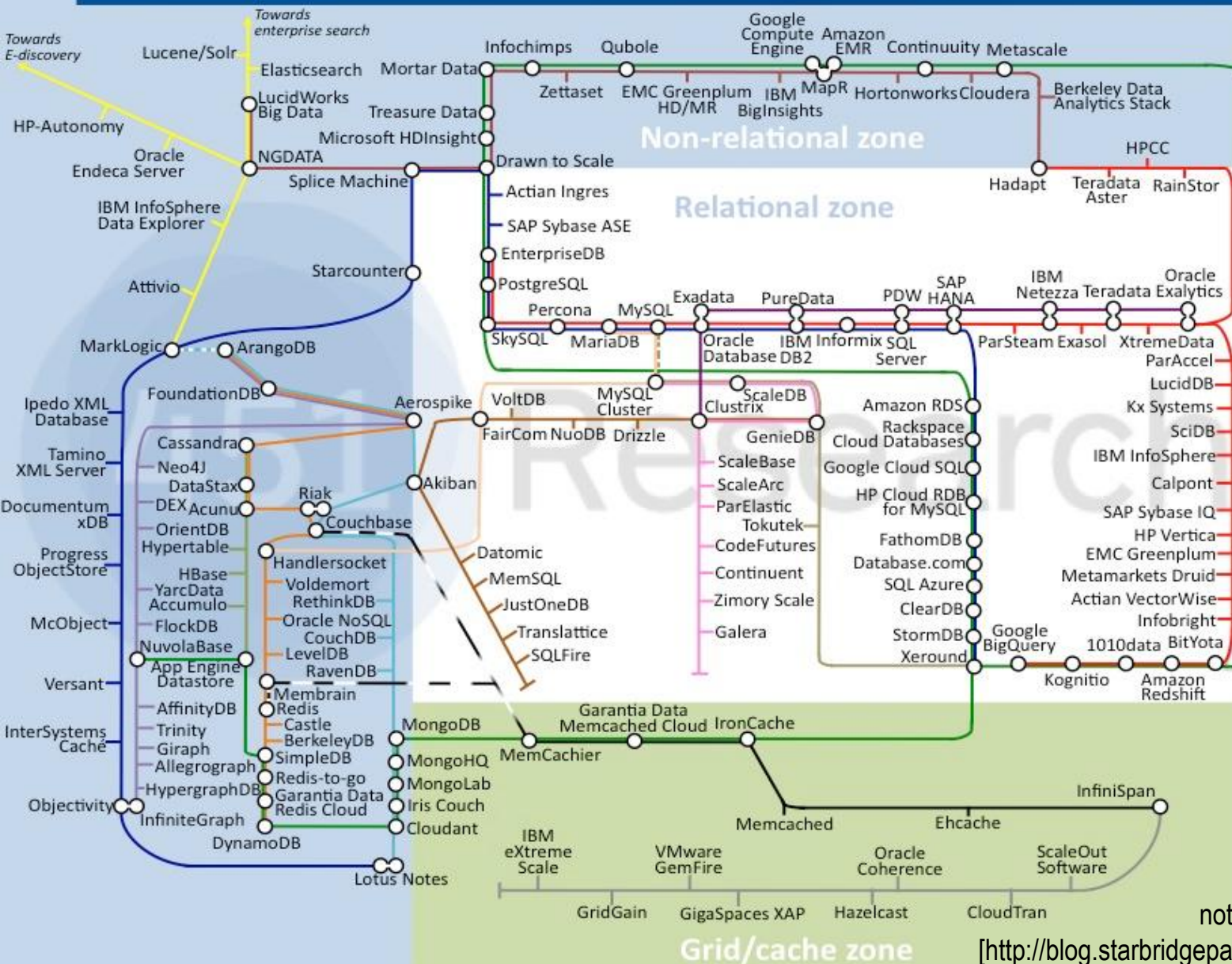
451 Research

Key:

- Operational
- Analytic
- as-a-Service
- - - NoSQL extension
- BigTables
- Graph
- Document
- Key value stores
- Key value direct access
- Hadoop
- - - NewSQL extension
- Storage engines
- Advanced clustering/sharding
- New SQL databases
- - - Data caching extension
- Data caching
- Data grid
- Index-based data management
- Appliances

www.451research.com

@maslett



not entirely correct/complete

[http://blog.starbridgepartners.com, 2013-aug19]

Summary & Outlook

- Fresh approach to scalable data services: NoSQL, NewSQL
 - Diversity of technology → pick best of breed for specific problem
- Avenue 1: **Modular data frameworks** to coexist
 - Heterogeneous model coupling barely understood - needs research
- Avenue 2: concepts **assimilated by relational vendors**
 - Like fulltext, object-oriented, SPARQL, ... cf „Oracle NoSQL“
- “SQL-as-a-service”
 - Amazon RDS, Microsoft SQL Azure, Google Cloud SQL
- *More than ever, experts in data management needed !*
 - *Both IT engineers and data engineers*