

SQL

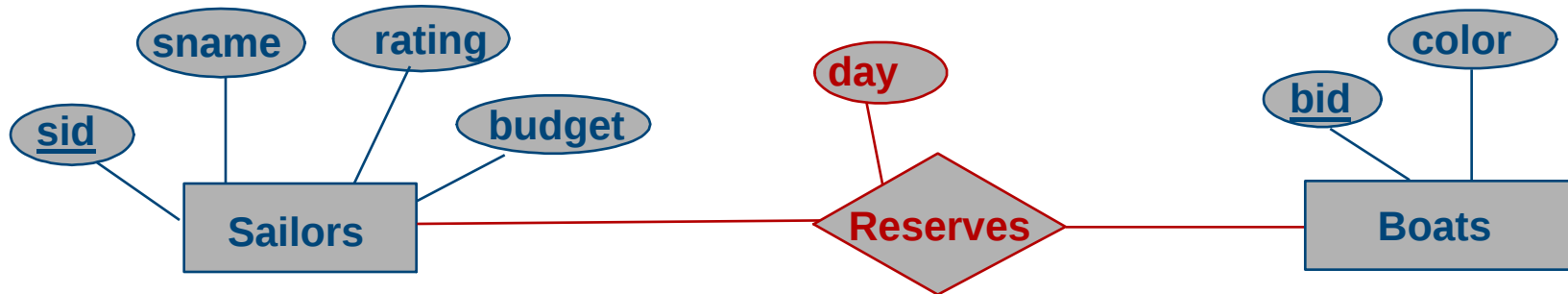
Ramakrishnan & Gehrke, Chapters 4 & 5



Sample Miniworld: Sailors & Boats



Example Instances



Sailors

sid	sname	rating	age
22	Dustin	7	45.0
31	Lubber	8	55.5
58	Rusty	10	35.0

Reserves

sid	bid	day
22	101	10/10/96
58	103	11/12/96

Boats

bid	color
101	red
102	blue
103	green

Basic SQL Query Structure

```
SELECT    [DISTINCT] target-list
FROM      relation-list
WHERE     qualification
```

- **relation-list**
 - list of relation names (possibly with a range-variable after each name)
- **target-list**
 - A list of attributes of relations in relation-list, possibly using **range variables**
- **qualification**
 - *Attr op const* or *Attr1 op Attr2* where op one of $<$, $>$, $=$, $\neq$, $\leq$, $\geq$ combined using AND, OR, NOT
- DISTINCT is optional for suppressing duplicates
 - By default duplicates not eliminated! ...so tables actually are **multisets**, not sets

Conceptual Evaluation Strategy

```
SELECT    [DISTINCT] target-list  
FROM      relation-list  
WHERE     qualification
```

- Semantics of an SQL query defined in terms of the following **conceptual evaluation strategy**:
 - Compute the **cross-product** of relation-list
 - Discard resulting tuples if they **fail qualification**
 - **Delete attributes** that are not in target-list
 - If DISTINCT is specified, **eliminate duplicate** rows
- This strategy is probably the **least efficient way** to compute a query!
 - An optimizer will find more efficient strategies to compute the same answers

Example of Conceptual Evaluation

↔ SELECT S.sname
FROM Sailors S, Reserves R
WHERE S.sid=R.sid AND R.bid=103

(sid)	sname	rating	age	(sid)	bid	day
22	Dustin	7	45.0	22	101	10/10/96
22	Dustin	7	45.0	58	103	11/12/96
31	Lubber	8	55.5	22	101	10/10/96
31	Lubber	8	55.5	58	103	11/12/96
58	Rusty	10	35.0	22	101	10/10/96
58	Rusty	10	35.0	58	103	11/12/96

■ cardinality?

A Note on Range Variables

- Really needed only if the **same relation appears twice** in the FROM clause
- previous query can also be written as:

```
SELECT S.sname  
FROM   Sailors S, Reserves R  
WHERE  S.sid=R.sid AND bid=103
```

- Or:

```
SELECT sname  
FROM   Sailors, Reserves  
WHERE  Sailors.sid=Reserves.sid AND bid=103
```

*It is good style,
however, to use
range variables
always!*

Join

- **Join** = several tables addressed in one query

```
SELECT target-list  
FROM Relation1 R1, Relation2 R2, ...  
WHERE qualification
```

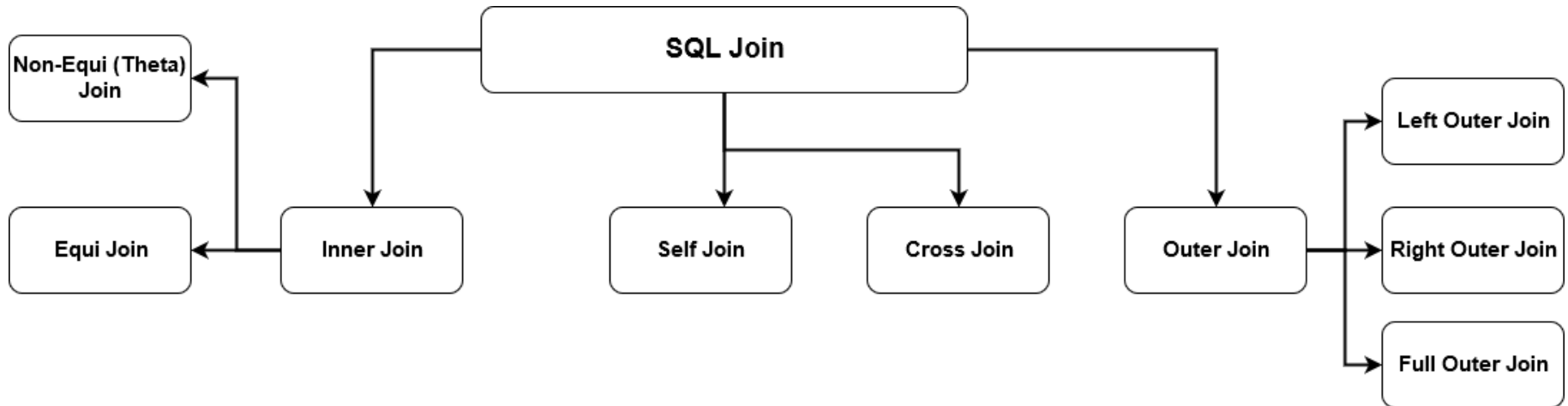
- List of relations in FROM clause determine **cross product**
- Frequently cross-relation **conditions** on attribute values to restrict results
- Most common: $R1.attr1 = R2.attr2$

- ex:

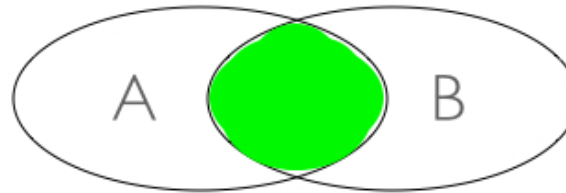
```
SELECT S.sid  
FROM Sailors S, Reserves R  
WHERE S.sid=R.sid
```

More Joins

- $T = R \bowtie_C S$
 - First build $R \times S$, then apply σ_C
- Generalization of equi-join: $A \theta B$ where θ one of $=, <, \dots$
 - Today, more general: σ_C can be any predicate
- Common join types [Quest]:



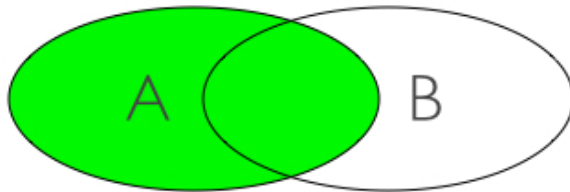
Even More on Joins



`SELECT * FROM A JOIN B ON A.id=B.id;`
`SELECT * FROM A, B WHERE A.id=B.id;`

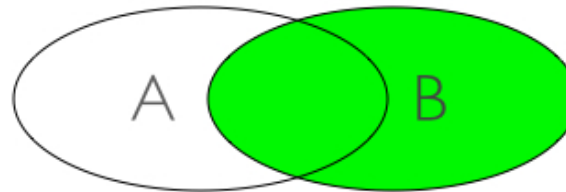
INNER JOIN

OUTER JOINS



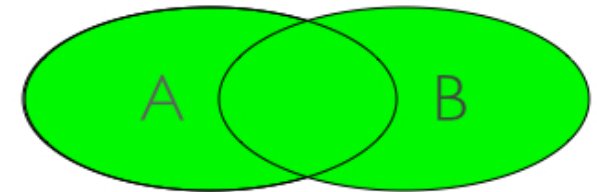
LEFT JOIN

`SELECT * FROM A LEFT JOIN B`
`ON A.id=B.id`



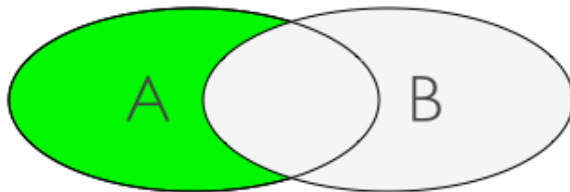
RIGHT JOIN

`SELECT * FROM A RIGHT JOIN B`
`ON B.id=A.id`

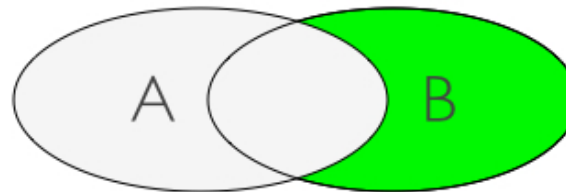


FULL JOIN

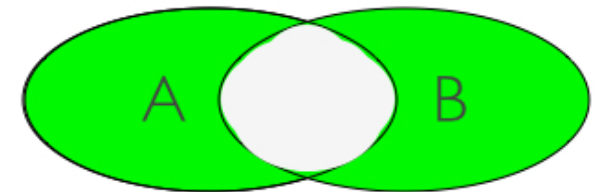
`SELECT * FROM A FULL JOIN B`
`ON A.id=B.id`



WHERE B.id IS NULL



WHERE A.id IS NULL



WHERE A.id IS NULL OR B.id IS NULL

"Sailors who've reserved at least 1 boat"

```
SELECT S.sid
FROM Sailors S, Reserves R
WHERE S.sid=R.sid
```

(sid)	sname	rating	age	(sid)	bid	day
22	Dustin	7	45.0	22	101	10/10/96
22	Dustin	7	45.0	58	103	11/12/96
31	Lubber	8	55.5	22	101	10/10/96
31	Lubber	8	55.5	58	103	11/12/96
58	Rusty	10	35.0	22	101	10/10/96
58	Rusty	10	35.0	58	103	11/12/96

- Would adding **DISTINCT** to this query make a difference?
- What is the effect of replacing **S.sid** by **S.sname** in the SELECT clause?
Would adding **DISTINCT** to this variant of the query make a difference?

"sid's of sailors who have reserved a red or a green boat"

- **UNION**: Can be used to compute the union of any two **union-compatible sets** of tuples
 - which themselves are the result of SQL queries
- If we replace **OR** by **AND** in the first version, what do we get?
- Also available: **EXCEPT**
 - What do we get if we replace UNION by EXCEPT?

```
SELECT S.sid
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid
      AND (B.color='red' OR B.color='green')
```

```
SELECT S.sid
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid
      AND B.color='red'
```

UNION

```
SELECT S.sid
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid
      AND B.color='green'
```

"Find sid's of sailors who have reserved a red and a green boat"

- **INTERSECT**: Can be used to compute the intersection of any two **union-compatible** sets of tuples
- Included in the SQL/92 standard, but some systems don't support it
- Contrast **symmetry** of the UNION and INTERSECT queries with how much the other versions differ!

```
SELECT S.sid
FROM Sailors S, Boats B1, Reserves R1,
      Boats B2, Reserves R2
WHERE S.sid=R1.sid AND R1.bid=B1.bid
      AND S.sid=R2.sid AND R2.bid=B2.bid
      AND (B1.color='red' AND B2.color='green')
```

Key field!

```
SELECT S.sid
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid
      AND B.color='red'

INTERSECT
SELECT S.sid
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid
      AND B.color='green'
```

Set Operations: Summary

- **SELECT** S1.a, S2.b FROM S1, S2
 - $S1 \times S2 = [\langle a, b \rangle \mid a \in S1, b \in S2]$
- S1 **UNION** S2
 - $S1 \cup S2 = [t \mid t \in S1 \vee t \in S2]$
- S1 **INTERSECT** S2
 - $S1 \cap S2 = [t \mid t \in S1 \wedge t \in S2]$
- S1 **EXCEPT** S2
 - $S1 \setminus S2 = [t \mid t \in S1 \wedge t \notin S2]$
- **SUM**(S.num), **AVG**(), ...
 - $\sum_{t \in S} t.num$
- **EXISTS**(S)
 - $S \neq \{ \}$
- t **IN** S2 t = ANY(S2)
 - $t \in S2$
- t op **ANY**(S) t op SOME(S)
 - $\exists x \in S: t \text{ op } x$
 - $(t \text{ op } s_1) \vee \dots \vee (t \text{ op } s_n)$ for $s_i \in S$
- t op **ALL** (S)
 - $\forall x \in S: t \text{ op } x$
 - $(t \text{ op } s_1) \wedge \dots \wedge (t \text{ op } s_n)$ for $s_i \in S$

Set Operations: Unique or Duplicates?

- Recall: Relations are multi-sets
- When are duplicates kept / eliminated?

keep duplicates	remove duplicates
SELECT	SELECT DISTINCT
UNION ALL	UNION
INTERSECT ALL	INTERSECT
EXCEPT ALL	EXCEPT

Aggregate Operators

- Summary information instead of value list

COUNT(*)
COUNT([DISTINCT] A)
SUM([DISTINCT] A)
AVG([DISTINCT] A)
MAX(A)
MIN(A)

A: single column

```
SELECT COUNT (*)
FROM Sailors S
```

```
SELECT COUNT (DISTINCT S.rating)
FROM Sailors S
WHERE S.sname='Bob'
```

```
SELECT AVG (S.age)
FROM Sailors S
WHERE S.rating=10
```

```
SELECT AVG ( DISTINCT S.age)
FROM Sailors S
WHERE S.rating=10
```

```
SELECT S.sname
FROM Sailors S
WHERE S.rating= (SELECT MAX(S2.rating)
FROM Sailors S2)
```

Nested Queries

- “Names of sailors who have reserved boat #103”:

```
SELECT S.sname
FROM Sailors S
WHERE S.sid IN (SELECT R.sid
                FROM Reserves R
                WHERE R.bid=103)
```

```
SELECT S.sname
FROM Sailors S, (SELECT R.sid
                 FROM Reserves R
                 WHERE R.bid=103) as X
WHERE S.sid = X.sid
```

- WHERE clause can itself contain an SQL query
 - so can FROM and HAVING, just not SELECT
- Sailors who have not reserved #103: use **NOT IN**
- To understand semantics of nested queries, think of a **nested loops evaluation**

“Sailors whose rating is greater than that of sailor Horatio”

- Nested:

```
SELECT *  
FROM Sailors S  
WHERE S.rating > ANY (SELECT S2.rating  
                      FROM Sailors S2  
                      WHERE S2.sname = 'Horatio')
```

- “flat” query:

```
SELECT S1.sname  
FROM Sailors S1, Sailors S2  
WHERE S1.rating > S2.rating and S2.sname = 'Horatio'
```

- Not every query can be flattened!

Unnesting Nested Queries

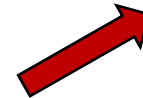
- Can sequentialize **uncorrelated** queries:

```
SELECT S.sname
FROM Sailors S
WHERE S.sid IN (SELECT R.sid
                FROM Reserves R
                WHERE R.bid=103)
```



- Always possible in FROM clause

```
SELECT S.sname
FROM Sailors S, (SELECT R.sid
                 FROM Reserves R
                 WHERE R.bid=103) as X
WHERE S.sid = X.sid
```



Common Table
Expression,
CTE (more)

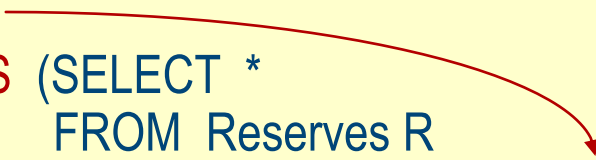
```
WITH Tmp AS (
  SELECT R.sid
  FROM Reserves R
  WHERE R.bid=103
)
SELECT S.sname
FROM Sailors S, Tmp
WHERE S.sid = Tmp.sid
```

- Not** in correlated subqueries ↩

Nested Queries with Correlation

- “Names of sailors who have reserved boat #103”:

```
SELECT S.sname
FROM Sailors S
WHERE EXISTS (SELECT *
               FROM Reserves R
               WHERE R.bid=103 AND S.sid=R.sid)
```



- **EXISTS**: another set operator, like IN
- If UNIQUE is used, and * is replaced by R.bid:
finds sailors with **at most one** reservation for boat #103
 - Why do we have to replace * by R.bid?
- Illustrates why, in general, subquery must be re-computed for each Sailors tuple

Nested Queries: INTERSECT vs IN

- "sailors who have reserved both red & green boat":

```
SELECT S.sid
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid AND B.color='red'
      AND S.sid IN (SELECT S2.sid
                    FROM Sailors S2, Boats B2, Reserves R2
                    WHERE S2.sid=R2.sid AND R2.bid=B2.bid
                      AND B2.color='green')
```

```
SELECT S.sid
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid
      AND B.color='red'

INTERSECT

SELECT S.sid
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid
      AND B.color='green'
```

- Similarly, EXCEPT queries re-written using NOT IN / NOT EXISTS ↩

Nested Queries: EXCEPT vs NOT EXISTS

- "sailors who have reserved all boats":
- Let's do it the hard way, without EXCEPT:

```
(1) SELECT S.sname
      FROM Sailors S
      WHERE NOT EXISTS
            ( (SELECT B.bid
              FROM Boats B)
              EXCEPT
              ( SELECT R.bid
                FROM Reserves R
                WHERE R.sid=S.sid ) )
```

```
(2) SELECT S.sname
      FROM Sailors S
      WHERE NOT EXISTS (SELECT B.bid
                        FROM Boats B
                        WHERE NOT EXISTS (SELECT R.bid
                                          FROM Reserves R
                                          WHERE R.bid=B.bid
                                          AND R.sid=S.sid ) )
```

Sailors S such that ...

there is no boat B without ...

a Reserves tuple showing S reserved B

More on Set Operators

- We have already seen IN, EXISTS and UNIQUE
 - Can also use NOT IN, NOT EXISTS and NOT UNIQUE
- Also available: *op ANY*, *op ALL*, *op* one of <, >, =, ≠, ≤, ≥
- "sailors whose rating is greater than that of sailor Horatio"

```
SELECT *  
FROM Sailors S  
WHERE S.rating > ANY (SELECT S2.rating  
                      FROM Sailors S2  
                      WHERE S2.sname = 'Horatio')
```

Expressions and Strings

```
SELECT S.age, age1=S.age-5, 2*S.age AS age2  
FROM Sailors S  
WHERE S.sname LIKE 'B_%B'
```

- Illustrates use of arithmetic expressions and string pattern matching:
 - Find triples (of ages of sailors and two fields defined by expressions) for sailors whose names begin and end with B and contain at least three characters
- AS and = are two ways to name fields in result
- LIKE is used for string matching
 - `'_'` stands for any one character
 - `'%'` stands for 0 or more arbitrary characters

Breaking the Set: ORDER BY

- So far: Query results are (multi) sets, hence **unordered**
Sometimes: need result **sorted**
- ORDER BY clause does this:

```
SELECT    [DISTINCT] target-list  
FROM      relation-list  
WHERE     qualification  
ORDER BY  sort-list [ASC|DESC]
```

- *sort-list*: list of attributes for ordering (ascending or descending order)
- Ex: “Names of all sailors,
in alphabetical order”

```
SELECT S.sname  
FROM   Sailors S  
ORDER BY S.sname
```

Grouping

- So far: aggregate operators applied to **all** (qualifying) tuples.
Sometimes: apply to each of several **groups** of tuples
- Consider: "age of the youngest sailor for each rating level"
 - Unknown # of rating levels, and rating values for levels
 - If we knew rating values go from 1 to 10:
can write loop of 10 queries:

```
For i = 1, 2, ... , 10:  
    SELECT MIN (S.age)  
    FROM Sailors S  
    WHERE S.rating = i
```

...or use **GROUP BY**:

```
SELECT      MIN( S.age )  
FROM        Sailors S  
GROUP BY    S.rating
```

Queries With GROUP BY and HAVING

```
SELECT    [DISTINCT] target-list
FROM      relation-list
WHERE     qualification
GROUP BY  grouping-list
HAVING    group-qualification
```

- *target-list* contains (i) attribute names, (ii) aggregate terms (ex: MIN(S.age))
- *grouping-list*: list of attributes for grouping
- *group-qualification*: group selection criterion (predicate on *grouping-list*)
- *target-list* attributes must be **subset of *grouping-list***
 - A **group** is a set of tuples that have the **same value for all attributes in *grouping-list***
 - Intuitively, each answer tuple corresponds to a group, and these attributes must have a single value per group

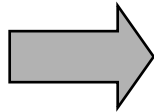
"Age of the youngest sailor with age ≥ 18 ,
for each rating with at least 2 such sailors"

```
SELECT S.rating, MIN (S.age) AS minage  
FROM Sailors S  
WHERE S.age >= 18  
GROUP BY S.rating  
HAVING COUNT (*) > 1
```

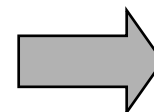
<u>sid</u>	sname	rating	age
22	dustin	7	45.0
29	brutus	1	33.0
31	lubber	8	55.5
32	andy	8	25.5
58	rusty	10	35.0
64	horatio	7	35.0
71	zorba	10	16.0
74	horatio	9	35.0
85	art	3	25.5
95	bob	3	63.5
96	frodo	3	25.5

"Age of the youngest sailor with age ≥ 18 ,
for each rating with at least 2 such sailors"

rating	age
7	45.0
1	33.0
8	55.5
8	25.5
10	35.0
7	35.0
10	16.0
9	35.0
3	25.5
3	63.5
3	25.5



rating	age
1	33.0
3	25.5
3	63.5
3	25.5
7	45.0
7	35.0
8	55.5
8	25.5
9	35.0
10	35.0



rating	minage
3	25.5
7	35.0
8	25.5

```
SELECT S.rating, MIN (S.age) AS minage
FROM Sailors S
WHERE S.age >= 18
GROUP BY S.rating
HAVING COUNT (*) > 1
```

Conceptual Evaluation

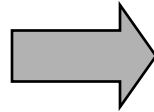
- compute cross-product of *relation-list*
- discard tuples that fail *qualification*
- delete `unnecessary' attributes
- **partition** remaining tuples into groups by value of attributes in *grouping-list*
- apply *group-qualification* to **eliminate** some **groups**
 - Expressions in *group-qualification* must have a single value per group!
- generate **one answer tuple** per qualifying group

SELECT	[DISTINCT] <i>target-list</i>
FROM	<i>relation-list</i>
WHERE	<i>qualification</i>
GROUP BY	<i>grouping-list</i>
HAVING	<i>group-qualification</i>

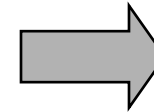
"Age of the youngest sailor with age ≥ 18 ,
for each rating with at least 2 such sailors
and with every sailor under 60"

HAVING COUNT (*) > 1
AND EVERY (S.age <=60)

rating	age
7	45.0
1	33.0
8	55.5
8	25.5
10	35.0
7	35.0
10	16.0
9	35.0
3	25.5
3	63.5
3	25.5



rating	age
1	33.0
3	25.5
3	63.5
3	25.5
7	45.0
7	35.0
8	55.5
8	25.5
9	35.0
10	35.0



rating	minage
7	35.0
8	25.5

What is the result of
changing EVERY
to ANY?

"Age of the youngest sailor with age ≥ 18 ,
for each rating with at least 2 sailors between 18 and 60"

```
SELECT S.rating, MIN (S.age) AS minage
FROM Sailors S
WHERE S.age >= 18 AND S.age <= 60
GROUP BY S.rating
HAVING COUNT (*) > 1
```

Answer relation:

rating	minage
3	25.5
7	35.0
8	25.5

Sailors instance:

<u>sid</u>	sname	rating	age
22	dustin	7	45.0
29	brutus	1	33.0
31	lubber	8	55.5
32	andy	8	25.5
58	rusty	10	35.0
64	horatio	7	35.0
71	zorba	10	16.0
74	horatio	9	35.0
85	art	3	25.5
95	bob	3	63.5
96	frodo	3	25.5

Quiz

HAVING does the same for GROUP BY groups
as _____ does for SELECT items

"For each red boat, the number of reservations for this boat"

```
SELECT B.bid, COUNT (*) AS scout
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid AND B.color='red'
GROUP BY B.bid
```

- Grouping over a join of three relations
- What if we remove B.color='red' from the WHERE clause and add a HAVING clause with this condition?
- What if we drop Sailors and the condition involving S.sid?

skipping

```
SELECT B.bid, COUNT (*) AS scout
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid
GROUP BY B.bid
HAVING B.color='red'
```

```
SELECT B.bid, COUNT (*) AS scout
FROM Boats B, Reserves R
WHERE R.bid=B.bid AND B.color='red'
GROUP BY B.bid
```

"Age of the youngest sailor with age ≥ 18 ,
for each rating with at least 2 sailors (of any age)"

```
SELECT S.rating, MIN(S.age)
FROM Sailors S
WHERE S.age >= 18
GROUP BY S.rating
HAVING (SELECT COUNT (*)
        FROM Sailors S2
        WHERE S.rating=S2.rating) > 1
```

- Shows HAVING clause can also contain a subquery
- Compare this with the query where we considered only ratings with 2 sailors over 18:
What if HAVING clause is replaced by:
 - HAVING COUNT(*) > 1

```
SELECT S.rating, MIN(S.age)
FROM Sailors S
WHERE S.age >= 18
GROUP BY S.rating
HAVING COUNT (*) > 1
```

"Those ratings for which the average age is the minimum over all ratings"

- Aggregate operations cannot be nested!

WRONG:

```
SELECT S.rating
FROM Sailors S
WHERE S.age = (SELECT MIN (AVG (S2.age))
               FROM Sailors S2)
```

- Correct solution (in SQL/92):

```
SELECT Temp.rating, Temp.avgage
FROM (SELECT S.rating, AVG (S.age) AS avgage
      FROM Sailors S
      GROUP BY S.rating) AS Temp
WHERE Temp.avgage = (SELECT MIN (Temp.avgage)
                    FROM Temp)
```

Null Values

- Field values in a tuple are sometimes **unknown** (e.g., a rating has not been assigned) or **inapplicable** (e.g., no spouse's name)
 - SQL provides a **special value null** for such situations
- Null **complicates** many issues, e.g.:
 - Special operators needed to check if value is/is not null
 - Is $\text{rating} > 8$ true or false when rating is equal to null?
 - *What about AND, OR and NOT connectives?*
 - We need a **3-valued logic** (true, false and unknown)
 - Meaning of constructs must be defined carefully
 - e.g., *WHERE* clause eliminates rows that don't evaluate to true
 - New operators (in particular, outer joins) possible/needed

Integrity Constraints (Review)

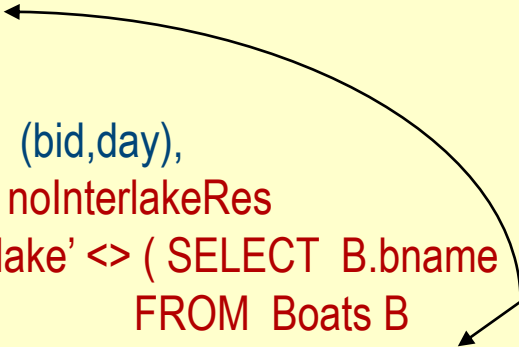
- IC describes conditions that every **legal instance** of a relation must satisfy
 - Inserts/deletes/updates violating ICs disallowed
 - Can be used to ensure application semantics (e.g., *sid* is a key), or prevent inconsistencies (e.g., *sname* has to be a string, *age* must be < 200)
- **Types of IC's**: Domain constraints, primary key constraints, foreign key constraints, general constraints
 - **Domain constraints**: Field values must be of right type. Always enforced

General Constraints

- Useful when more general ICs than keys are involved
- Can use queries to express constraint
- Constraints can be named

```
CREATE TABLE Sailors
(
  sid INTEGER,
  sname CHAR(10),
  rating INTEGER,
  age REAL,
  PRIMARY KEY (sid),
  CHECK ( rating >= 1 AND rating <= 10 )
)
```

```
CREATE TABLE Reserves
(
  sname CHAR(10),
  bid INTEGER,
  day DATE,
  PRIMARY KEY (bid,day),
  CONSTRAINT noInterlakeRes
  CHECK ('Interlake' <> ( SELECT B.bname
                        FROM Boats B
                        WHERE B.bid=bid) )
)
```



Assertions

- CHECK constraint is awkward and **wrong!**
- If Sailors is empty, number of Boats tuples can be *anything*
- **ASSERTION** is the right solution: not associated with either table

```
CREATE TABLE Sailors
(  sid INTEGER,
  sname CHAR(10),
  rating INTEGER,
  age REAL,
  PRIMARY KEY (sid),
  CHECK
  ( (SELECT COUNT (S.sid) FROM Sailors S)
    + (SELECT COUNT (B.bid) FROM Boats B) < 100 )
)
```

*Number of boats
+ number of sailors
is < 100*

```
CREATE ASSERTION smallClub
CHECK
( (SELECT COUNT (S.sid) FROM Sailors S)
  + (SELECT COUNT (B.bid) FROM Boats B) < 100
)
```

Triggers

- **Trigger:** procedure that starts automatically if & when specified changes occur to the database
- Three parts ("ECA rules"):
 - **Event** -- activates the trigger
 - **Condition** -- tests whether the triggers should run
 - **Action** -- what happens if the trigger runs

Triggers: Example (SQL:1999)

```
CREATE TRIGGER youngSailorUpdate
  AFTER INSERT ON Sailors
  REFERENCING NEW TABLE NewSailors
  FOR EACH STATEMENT
  INSERT
    INTO YoungSailors( sid, name, age, rating )
  SELECT sid, name, age, rating
  FROM NewSailors N
  WHERE N.age <= 18
```

Summary

- SQL important factor for acceptance of relational model
 - more natural than earlier, procedural query languages
 - Simple, easy-to-grasp paradigm: sets + few generic operations on them
 - **Relationally complete** = as powerful as relational algebra
 - *in fact, significantly more expressive power than relational algebra*
 - Not *computationally* complete! (no recursion, for example)
- **Set orientation** good basis for declarative query language
 - **Declarative** = describe desired **result** (well, almost :-), more user-oriented (imperative = describe algorithm; more implementation-oriented)
- SQL allows specification of **integrity constraints**
- **Triggers** respond to changes in the database

Summary (Contd.)

- Many **alternative phrasings**
 - optimizer should look for most efficient evaluation plan
 - In practice, users need to be aware of how queries are optimized and evaluated for best results
- **NULL** for unknown field values
 - brings many complications
- ...and we have left out a lot!
 - Recursion, PL/SQL, schema evolution, ...